

An analytical approach for calculating end-to-end response times in autonomous driving applications

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Agenda

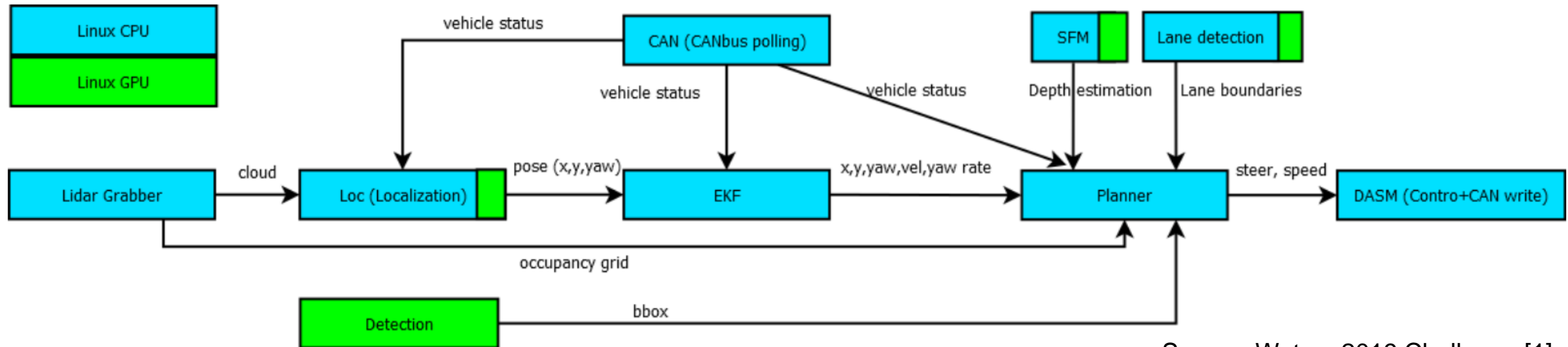
- Introduction
- Data consistency
- Data propagation paths
- Analysis Approach
- Optimization
- Integrated Analysis and Optimization Results
- Conclusion and outlook

Introduction

"Boosting Design Efficiency for Heterogeneous³ Systems"

Programcall	ITEA 3 Call 4 17003
Title	Boosting Design Efficiency for Heterogeneous ³ Systems
Period	Apr 2019 - Mar 2022
Status	Running
Domain	Services, Systems & Software Creation
Technology	Software
Effort	122 man-years
Costs	EUR 15.9 million
Project Leader	Jörg Tessmer (Bosch)
Partners	25
Countries (5)	Finland, Germany, Portugal, Sweden, Turkey https://itea3.org/project/panorama.html

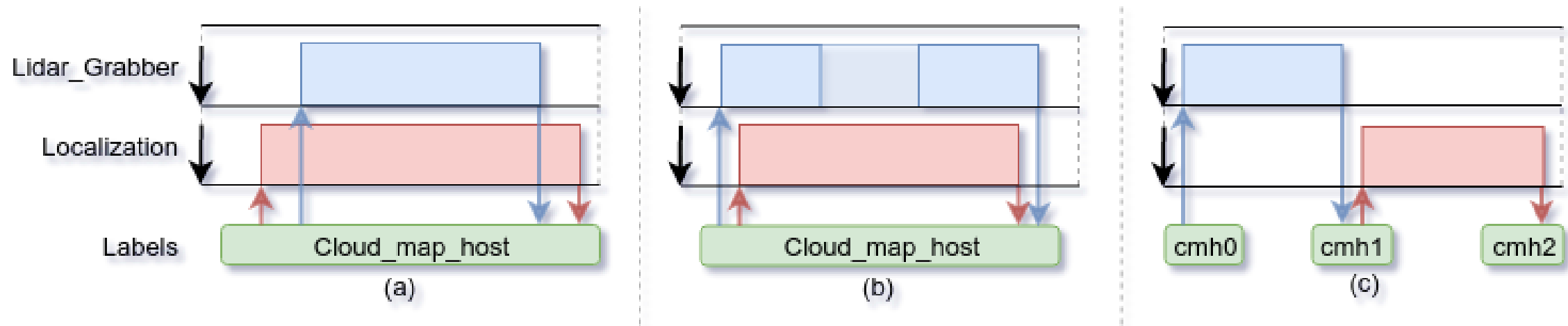




Source: Waters 2019 Challenge [1]

- Calculate applications end-to-end response time
 - Derive task chains for end-to-end paths
 - Develop integrated response time analysis approach
- Optimize the latency of the different task-chains
 - Our scope: Minimize the end-to-end response time

Data consistency



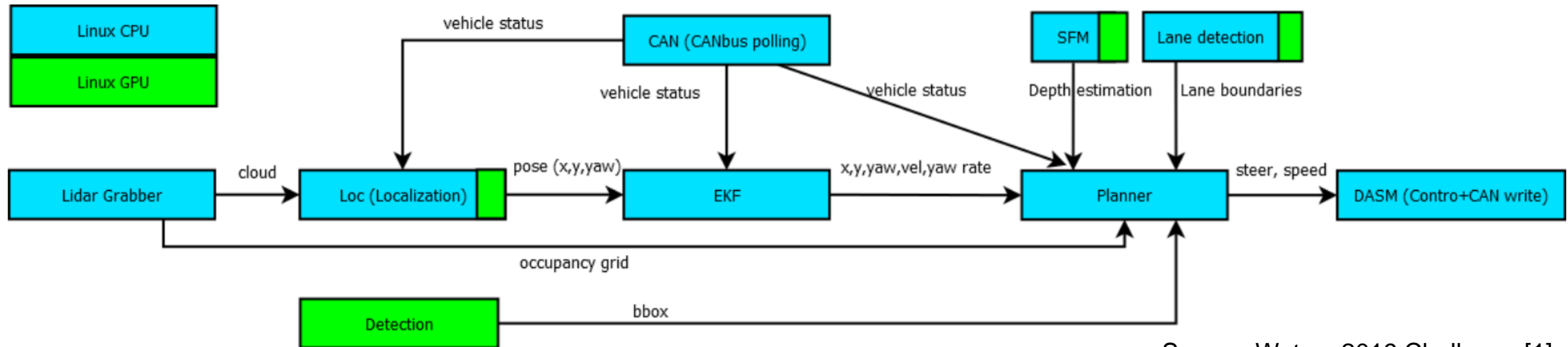
(a) Localization overwrites Lidar_Grabber

(b) Lidar_Grabber overwrites Localization

(c) Deterministic behaviour

- Higher memory consumption
- Increased latency compared to e.g. semaphore usage
- Correct behaviour can be realised at the cost of higher latency by an e.g. pipeline fashioned approach

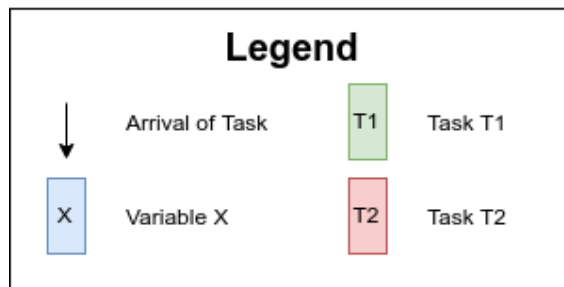
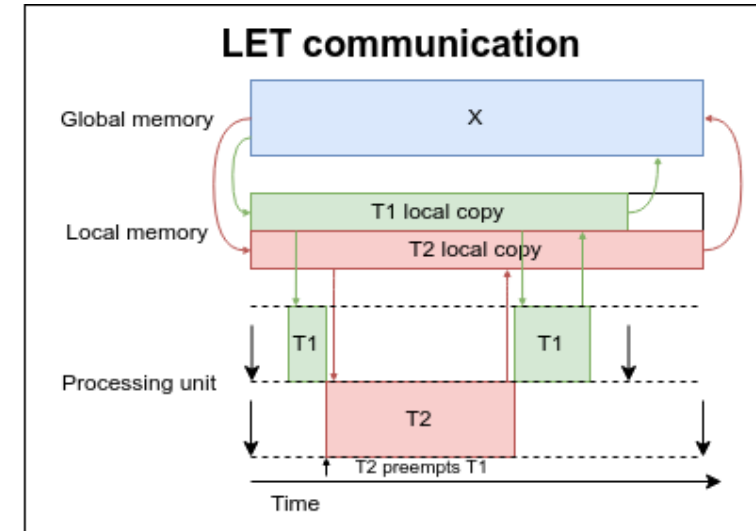
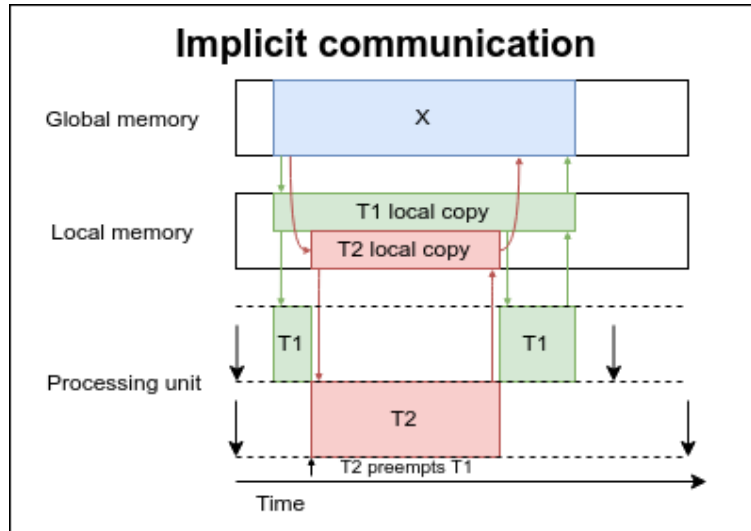
Data Propagation Paths



Source: Waters 2019 Challenge [1]

■ All critical paths from sensor tasks to actuator tasks

- Lidar_Grabber → Loc → EKF → Planner → DASM
- CAN → Loc → EKF → Planner → DASM
- SFM → Planner → DASM
- Lane_detection → Planner → DASM
- Detection → Planner → DASM

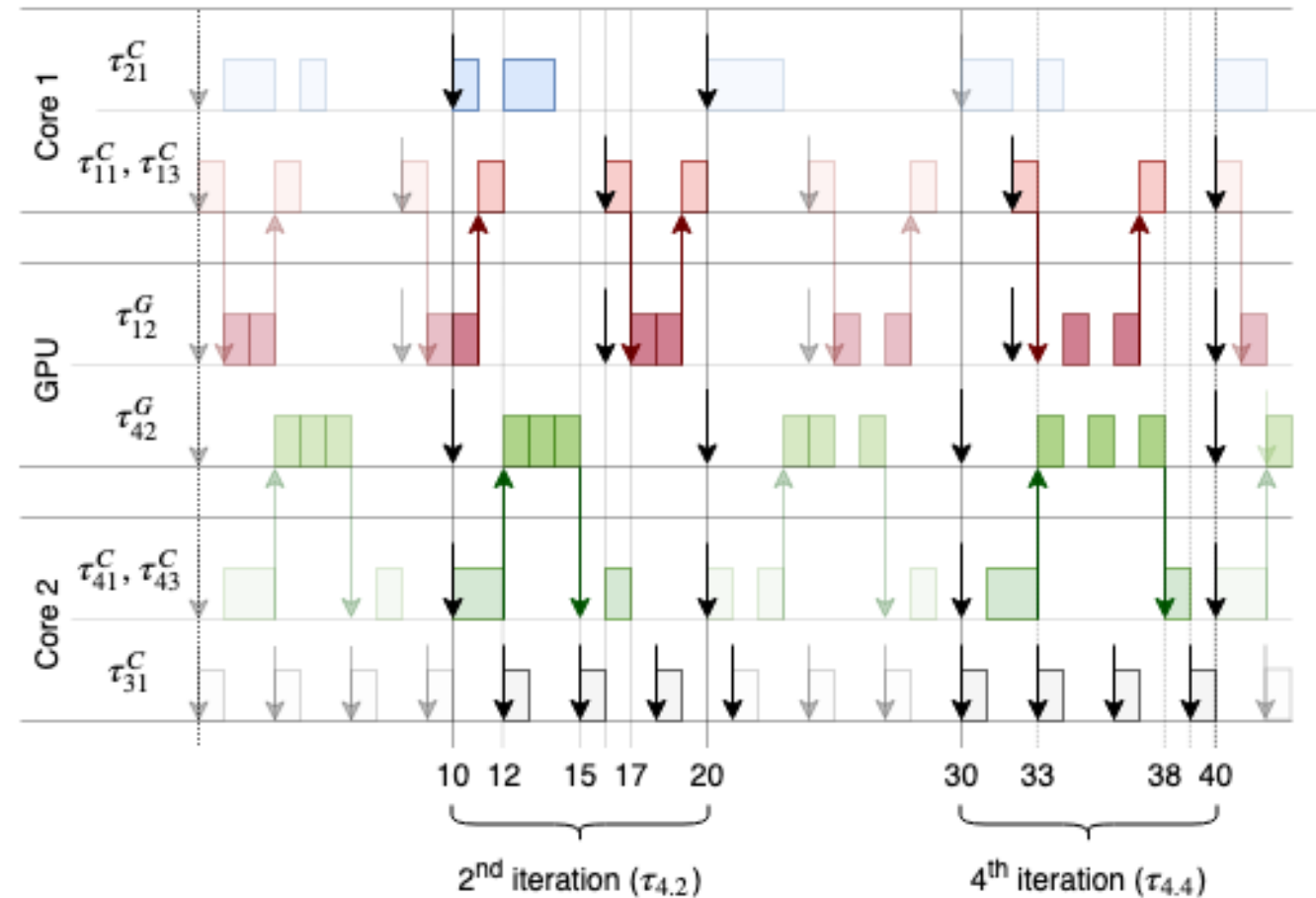


- Implicit communication
 - End-to-end response time can be optimized
 - E2E-RT by Kloda et al. [2]
- LET communication
 - Deterministic behaviour
 - Own implementation extending [2]

- Different scheduling strategies
 - Fixed priority preemptive (FPP) scheduling on CPUs
 - Weighted round-robin (WRR) scheduling on GPUs
 - Task suspension

- FPP: Palencia et al. [3]

- WRR: Racu et al. [4]



Analysis – Task Model

- Tasks are described in terms of **transactions**, with:
(Sub-Tasks (Runnables), Period, Priority)

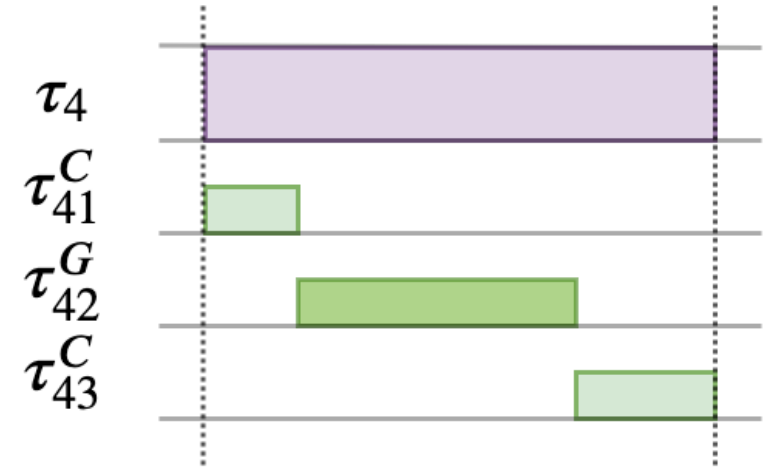
$$\tau_i = (\{\tau_{i1}, \dots, \tau_{i|\tau_i|}\}, P_i, \pi_i)$$

- Sub-Task on **CPU**, with:
(Execution Time, Offset, Jitter)

$$\tau_{ij}^C = (C_{ij,\rho}, O_{ij}, J_{ij})$$

- Sub-Task on **GPU**, with:
(Execution Time, Offset, Jitter, Time-Slice)

$$\tau_{ij}^G = (C_{ij,\rho}, O_{ij}, J_{ij}, \phi_{ij})$$



Analysis – Data transfer times

- Number of **label accesses**

$$\lambda_{ij} = \sum_{l \in \mathcal{L}_{ij}} \left\lceil \frac{\text{size}(l)}{\text{size}(\text{cacheline})} \right\rceil$$

- Memory **access times**

	Best	Worst
A57	20 ns	220 ns
Denver	8 ns	38 ns
GPU	3 ns	6 ns

- Total **work** for a sub-task

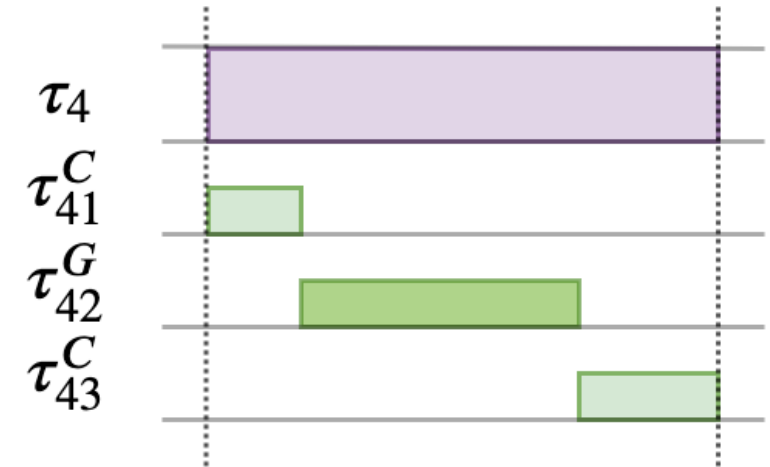
$$W_{ij} = C_{ij,\rho} + \lambda_{ij} \cdot A_{\rho}$$

- Sub-task's best case response time

$$\mathcal{R}_{ij}^+ = \sum_{k=1 \dots j} W_{ik}^+$$

- Task's worst case response time

$$\mathcal{R}_i^- = \mathcal{R}_{i|\tau_i}^-$$

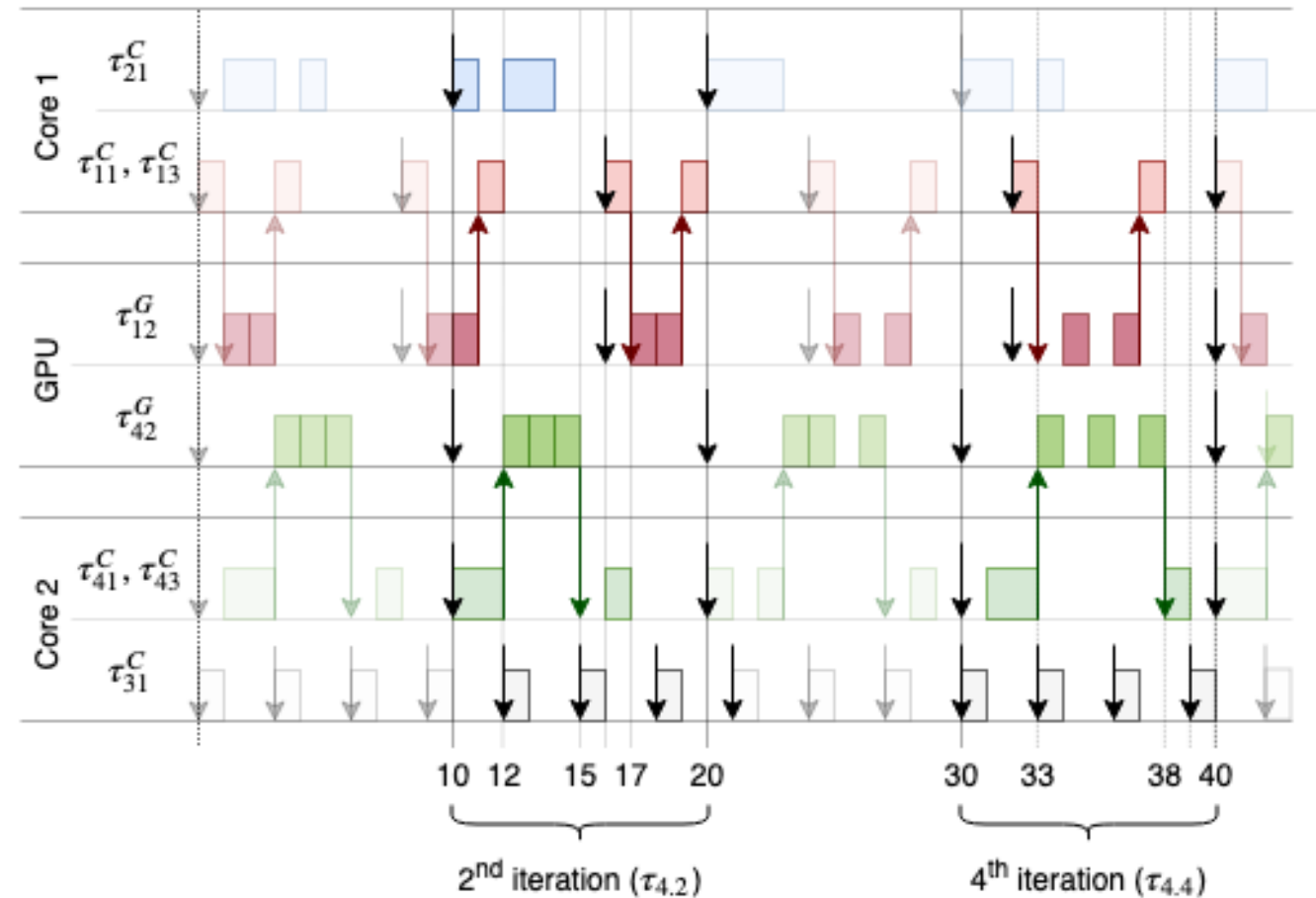


- Integration of both approaches as iterative strategy
 - Update the **offset** of the **successor**, set it to the **best case response time** of its predecessor

$$O_{ij} = R_{ij-1}^+$$

- Update the **jitter**, set it to the difference between **worst case response time** and **offset (BCRT)**

$$J_{ij} = R_{ij-1}^- - R_{ij-1}^+$$



Optimization

- Genetic Algorithm Implementation based on Jenetics (Java)
- Already integrated into App4MC (OpenMapping)
- Degrees of freedom (DoF)
 - Allocation (Task to Processing Unit)
 - Allocation (Offloadable sub-task to Processing Unit)
 - Time Slice (Sub-Task on GPU only)

Integrated Analysis and Optimization Results

- Similar end-to-end latency for LET and implicit communication
- Response times close to the task's period
- Runtime: 287 seconds
 - Reason: Audsleys priority assignment algorithm

Task Chain	LET end-to-end	Implicit end-to-end
σ_1	886	859.9
σ_2	865	836.9
σ_3	67	59.9
σ_4	100	71.9
σ_5	230	221.9

Name	P	π	C^-	$\lambda \cdot \mathcal{A}^-$	R^-	ϕ
Core 0 (Denver)						
Planner	12	9	11.2	0.8	12.0	—
Core 1 (Denver)						
SFM*	33	6	6.7	3.6	31.5	—
Lane_detection	66	2	42.2	1.2	53.6	—
Core 2 (A57)						
CANbus_polling	10	5	0.6	0.0	0.6	—
EKF	15	1	4.8	0	5.4	—
Core 3 (A57)						
Localization	400	4	387.4	5.2	392.6	—
Core 4 (A57)						
Lidar_Grabber	33	8	13.7	12.0	25.7	—
Detection*	200	7	4.7	1.8	198.0	—
Core 5 (A57)						
OS_Overhead	100	0	50	0.0	79.9	—
DASM	5	3	1.9	0.0	1.9	—
GP10B (GPU)						
Detection	200	—	116.0	0.5	170.5	7.0
SFM	33	—	7.9	0.4	15.2	11.6

Conclusion and outlook

- Analysis of end-to-end response time of a given application following an **implicit** and **LET** communication paradigm
- Accounting all mandatory delays:
 - Data transfer time for **copy engine** (GPU <-> CPU)
 - Data transfer time between **CPU and shared main memory**
 - **Synchronous** and **asynchronous** offloading
 - Application of given memory contention approach
- Response time analysis for coupled task sets scheduled on an **heterogeneous architecture** consisting of processing units with **fixed priority preemptive** (CPU) and **weighted round robin** (GPU) scheduling
- **Minimization** of the applications maximum **end-to-end response** time among all task chains for a implicit communication paradigm

Conclusion and outlook

- Simplification of the model was required (transitive labels, planner task)
- Cooperative scheduling (FPFP optimistic assumption)
- Scalability
- Fully integrated approach
- Comparison with practical demonstration results

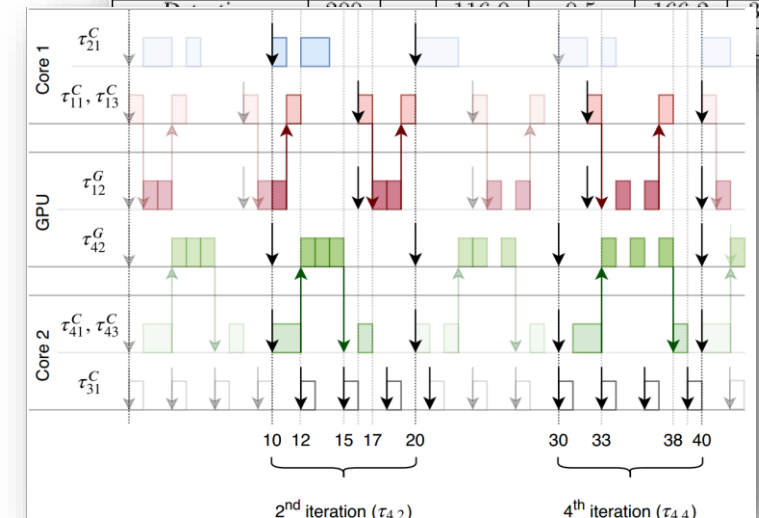
References

- [1] Arne Hamann, Dakshina Dasari, and Falk Wurst. WATERS Industrial Challenge, 2019.
- [2] Tomasz Kloda, Antoine Bertout, and Yves Sorel. Latency analysis for data chains of real-time periodic tasks. In 23rd IEEE International Conference on Emerging Technologies and Factory Automation, ETFA 2018, Torino, Italy, September 4-7, 2018, pages 360–367. IEEE, 2018.
- [3] J.C. Palencia and M. Gonzalez Harbour. Schedulability analysis for tasks with static and dynamic offsets. 2002.
- [4] Razvan Racu, Li Li, Rafik Henia, Arne Hamann, and Rolf Ernst. Improved response time analysis of tasks scheduled under preemptive round-robin. In Proceedings of the 5th International Conference on Hardware/Software Codesign and System Synthesis, CODES+ISSS 2007, Salzburg, Austria, September 30 - October 3, 2007, pages 179–184. ACM, 2007.

Thank you for your attention

- Analysis of End-to-End latencies of a given application following an **implicit** and **LET** communication paradigm
- Response time analysis for coupled task sets scheduled on an **heterogeneous architecture** consisting of processing units with **fixed priority preemptive** (CPU) and **weighted round robin** (GPU) scheduling
- **Minimization** of the applications maximum **end-to-end response time** among all task chains for a implicit communication paradigm

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Questions?

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